

Quantum Time-Space Tradeoffs for Exponential Dynamic Programming

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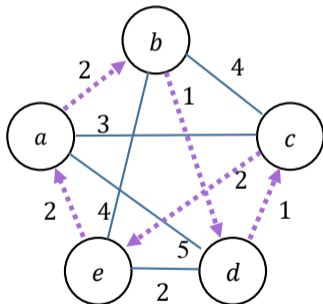
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Travelling Salesman Problem

- Given a weighted graph on n vertices, find the shortest cycle that passes through each vertex exactly once.



A Question (Theory Days 2018)

- TRAVELLING SALESMAN PROBLEM:

Classically: [Bel62, HK62]

$$\tilde{O}(2^n)$$

Quantumly: [ABI⁺19]

$$\tilde{O}(1.728^n)$$

- *“Can the classical time-space tradeoffs of Koivisto and Parviainen be applied here? [KP10]”*

— Someone in the audience.

Time-Space Tradeoffs

- Less space 😊, more time 😞.
- Why? Let D_1, \dots, D_N be N memory records.
- Quantum Random Access Memory (QRAM):

$$\sum_{i=1}^N \alpha_i |i\rangle_{\text{address}} |0\rangle_{\text{target}} \longmapsto \sum_{i=1}^N \alpha_i |i\rangle_{\text{address}} |D_i\rangle_{\text{target}}$$

- $O(\text{polylog } N)$ in theory [GLM08], very hard to implement! [JR25]

[KP10] Time-Space Tradeoffs

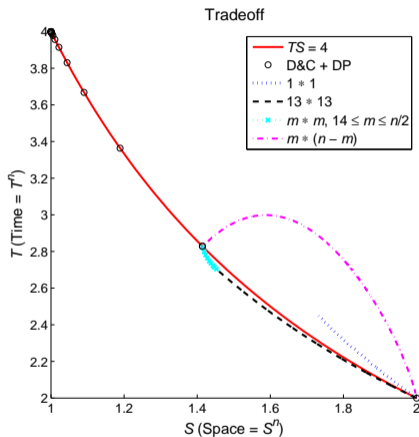


Figure 1: Space-time tradeoff schemes for permutation problems. The time requirement $O^*(T^n)$ is shown as a function of the space requirement $O^*(S^n)$, for $1 \leq S \leq 2$.

Roadmap

1. Classical Tradeoff \longrightarrow Quantum Algorithm.

Gurevich and Shelah Tradeoff

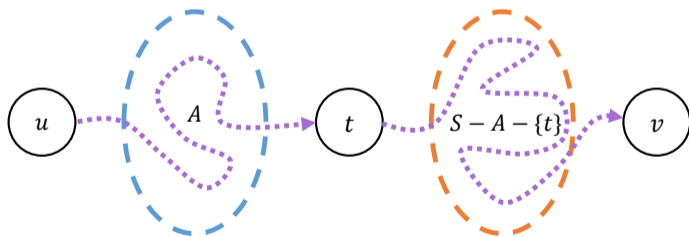
- TSP Algorithm 1: Divide & Conquer. Let $G = (V, E)$.
- $f(a, S, b)$ — the shortest path from a to b , visiting all of S .
- Optimal route:

$$\text{OPT} = \min_{\substack{e \in V \\ \{1, e\} \in E}} \left(f(1, V - \{1\} - \{e\}, e) + w(1, e) \right).$$

Gurevich and Shelah Tradeoff

- D&Q recurrence:

$$f(u, S, v) = \min_{\substack{ACS \\ |A|=|S|/2 \\ t \in S-A}} \left(f(u, A, t) + f(t, S - A - \{t\}, v) \right).$$



Gurevich and Shelah Tradeoff

- TSP Algorithm 1: Divide & Conquer. Let $G = (V, E)$.
- Time complexity:

$$\binom{n}{n/2} \binom{n/2}{n/4} \cdots \binom{2}{1} = O(2^n \cdot 2^{n/2} \cdots 2) = O(4^n).$$

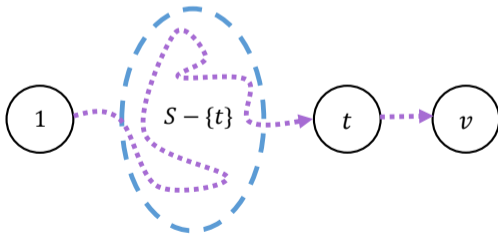
- Space complexity:

$$n + n/2 + \dots + 1 = O(n).$$

Gurevich and Shelah Tradeoff

- TSP Algorithm 2: Dynamic Programming.
- DP recurrence:

$$f(1, S, v) = \min_{t \in S} \left(f(1, S - \{t\}, t) + w(t, v) \right).$$



Gurevich and Shelah Tradeoff

- TSP Algorithm 2: Dynamic Programming.
- Time complexity:

$$|\{S \mid S \subseteq V\}| \cdot n^2 = \tilde{O}(2^n).$$

- Space complexity:

$$|\{S \mid S \subseteq V\}| \cdot n = \tilde{O}(2^n).$$

Gurevich and Shelah Tradeoff [GS87]

- TSP Tradeoff: pick a $k \in [\log_2 n]$.
 1. Divide & Conquer while $S > 2^k$.
 2. Dynamic Programming when $S = 2^k$.
- Time complexity: $\tilde{O}(4^n / 2^k)$.
- Space complexity: $\tilde{O}(2^k)$.

Quantum Algorithm?

- Grover's search: for x_1, \dots, x_N can find

$$\min_{i=1}^N x_i$$

in $\tilde{O}(\sqrt{N})$ time and $O(\log N)$ space.

Quantum Algorithm?

1. Dynamic Programming until sets of size αn for some $\alpha \in [0, 1/2]$.
2. Grover's search over the Divide & Conquer!
 - Time complexity:

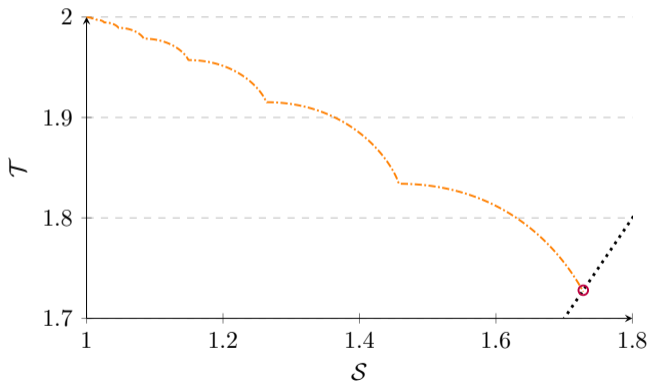
$$\tilde{O}\left(\binom{n}{\alpha n} + \sqrt{\binom{n}{n/2} \binom{n/2}{n/4} \cdots \binom{n/2^{\log(1/\alpha)}}{\alpha n}}\right).$$

- Space complexity (QRAM):

$$\tilde{O}\left(\binom{n}{\alpha n}\right).$$

Quantum Tradeoff

- By optimizing α , we get $\mathcal{T} = \mathcal{S} = \tilde{O}(1.728^n)$ [ABI⁺19].
- Other α choices give a time-space tradeoff.



Quantum Tradeoff

- Is this the best we can do?

Quantum Tradeoff

- Let's use Gurevich and Shelah again!

Roadmap

1. Classical Tradeoff \longrightarrow Quantum Tradeoff.
2. Quantum Tradeoff
+
Classical Tradeoff
(again) \longrightarrow (an improved)
Quantum Tradeoff.

Improved Quantum Tradeoff

The algorithm: pick two parameters $\alpha \in [0, 1/2]$ and $\beta \in [0, 1/2]$.

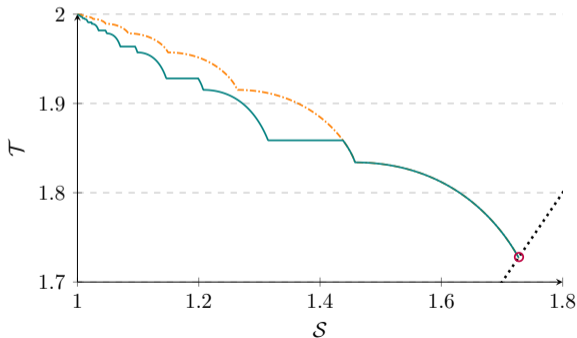
1. Run Grover's search + Divide & Conquer for sets larger than βn .
2. Else, use the [ABI⁺19] quantum algorithm with parameter α .

Improved Quantum Tradeoff

- By optimizing α and β , we obtain complexity

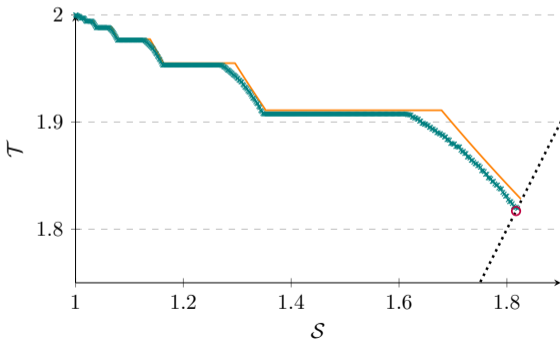
$$\mathcal{S} = \tilde{O}\left(2^{H(\alpha)n}\right), \quad \mathcal{T} = \tilde{O}\left(2^{\max\left\{1 - \frac{2 - H(2^k \alpha)}{2^{k+1}}, H(\alpha)\right\}n}\right).$$

- An improved tradeoff:



Quantum Tradeoff for Permutation Problems

- We also optimize the permutation problem quantum algorithm of [ABI⁺19], for PATHWIDTH, CUTWIDTH etc. (teal).
- We then “quantize” the *pairwise scheme* [PK09] time-space tradeoff (orange):



P.S. New Classical Improvements

- Koivisto and Parviainen: $(\mathcal{TS})^{1/n} \leq 3.93$ [KP10].
- arxiv.org/abs/2604.05661: $(\mathcal{TS})^{1/n} \leq 3.75$.
- arxiv.org/abs/2604.05645: $(\mathcal{TS})^{1/n} \leq 3.58$.

Thank you!
Questions?



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